

بسم الله الرحمن الرحيم

فصل پنجم

زبان‌های مستقل از متن (۲)

Context-Free Languages (2)

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دانشگاه تهران



Compilers

Program

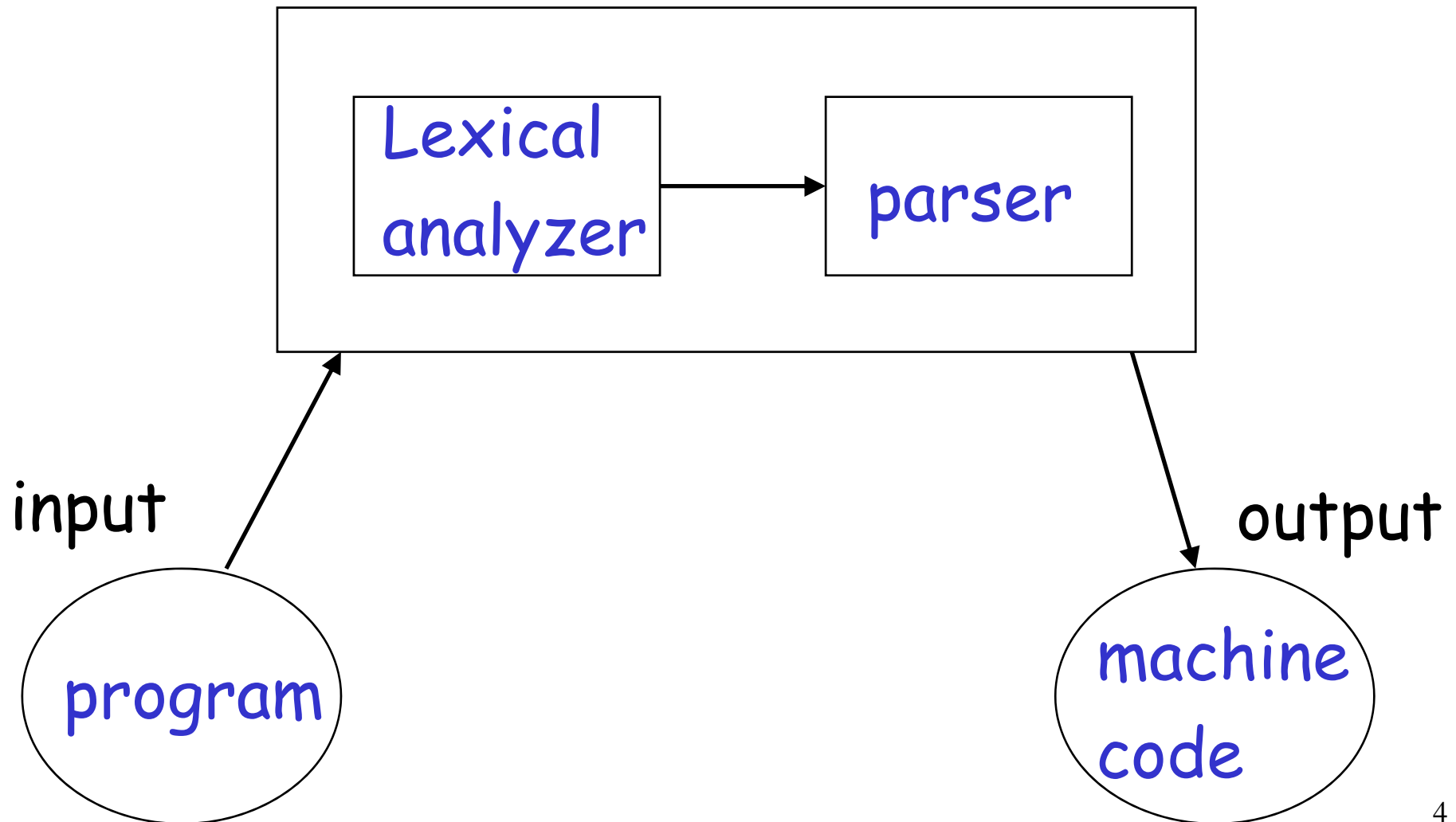
```
v = 5;  
if (v>5)  
    x = 12 + v;  
while (x !=3) {  
    x = x - 3;  
    v = 10;  
}  
.....
```

Compiler

Machine Code

```
Add v,v,0  
cmp v,5  
jmplt ELSE  
THEN:  
    add x, 12,v  
ELSE:  
    WHILE:  
        cmp x,3  
    ...
```

Compiler



A **parser** knows the grammar
of the programming language

Parser

$\text{PROGRAM} \rightarrow \text{STMT_LIST}$

$\text{STMT_LIST} \rightarrow \text{STMT}; \text{STMT_LIST} \mid \text{STMT};$

$\text{STMT} \rightarrow \text{EXPR} \mid \text{IF_STMT} \mid \text{WHILE_STMT}$
 $\mid \{ \text{STMT_LIST} \}$

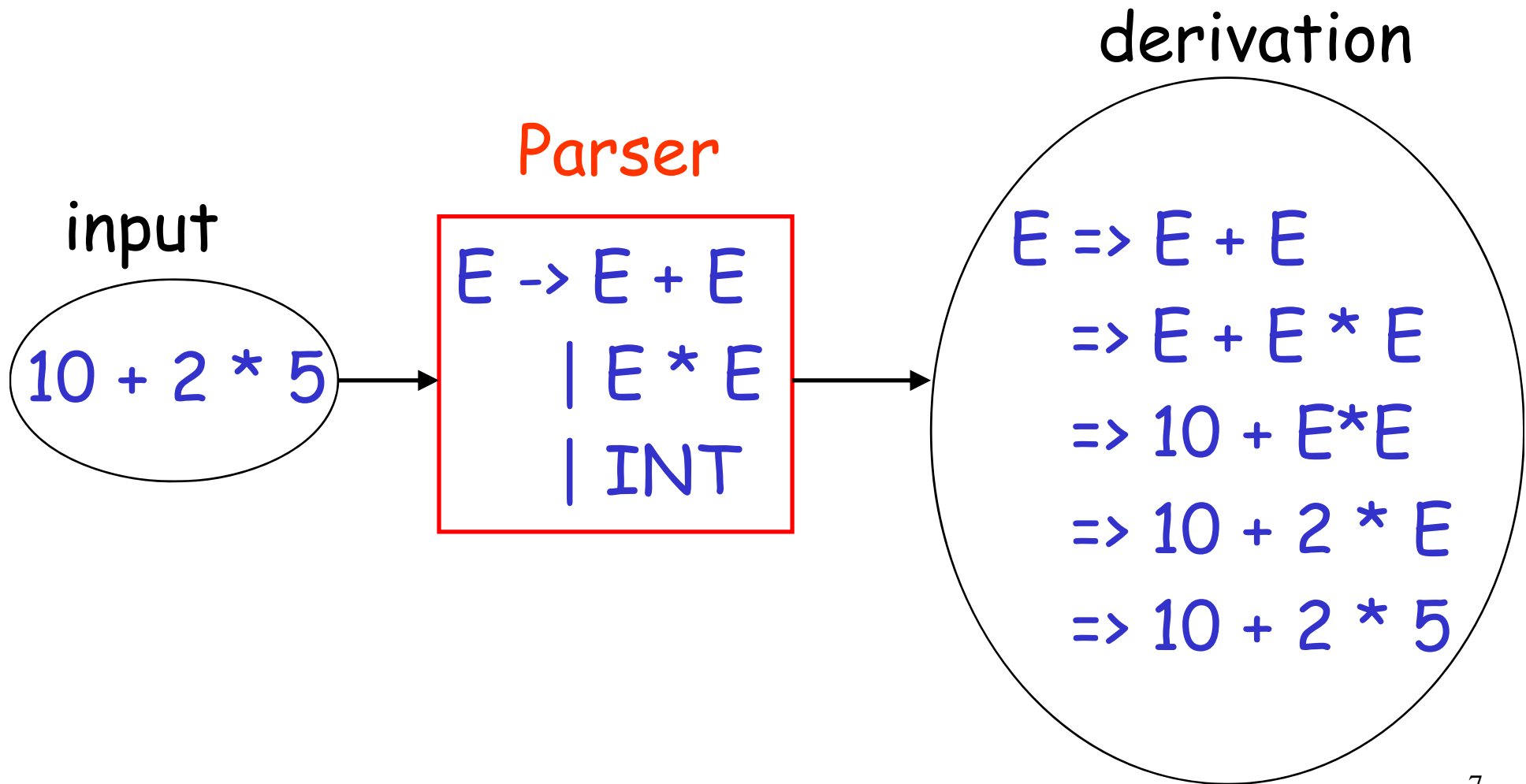
$\text{EXPR} \rightarrow \text{EXPR} + \text{EXPR} \mid \text{EXPR} - \text{EXPR} \mid \text{ID}$

$\text{IF_STMT} \rightarrow \text{if (EXPR) then STMT}$

$\mid \text{if (EXPR) then STMT else STMT}$

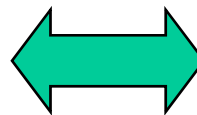
$\text{WHILE_STMT} \rightarrow \text{while (EXPR) do STMT}$

The parser finds the derivation
of a particular input

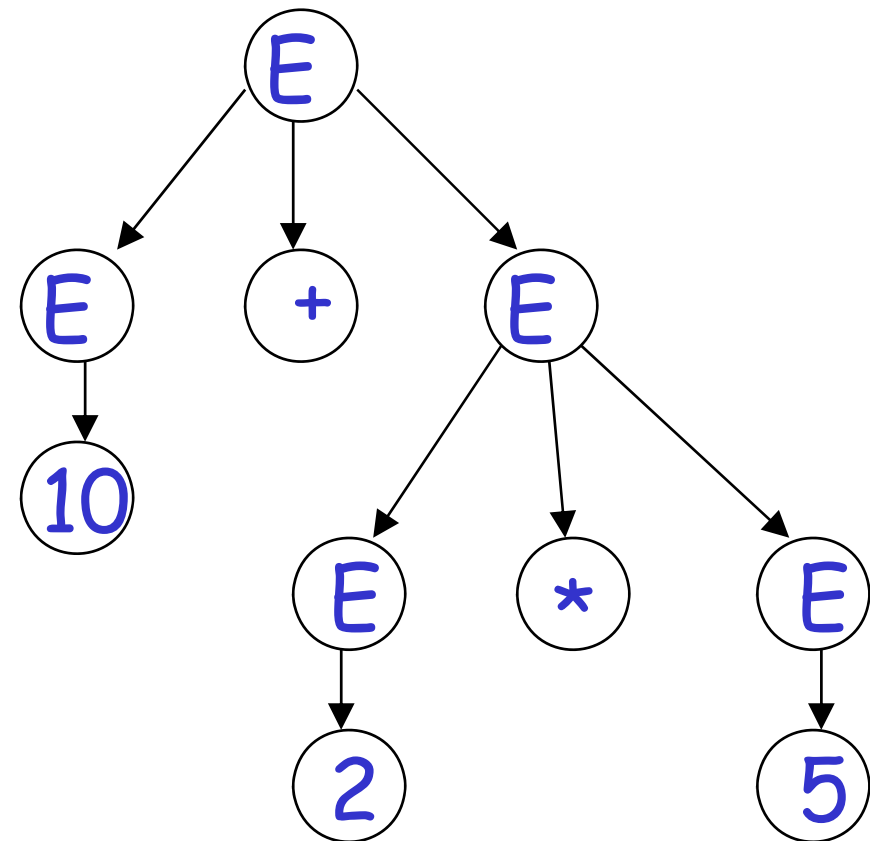


derivation

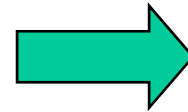
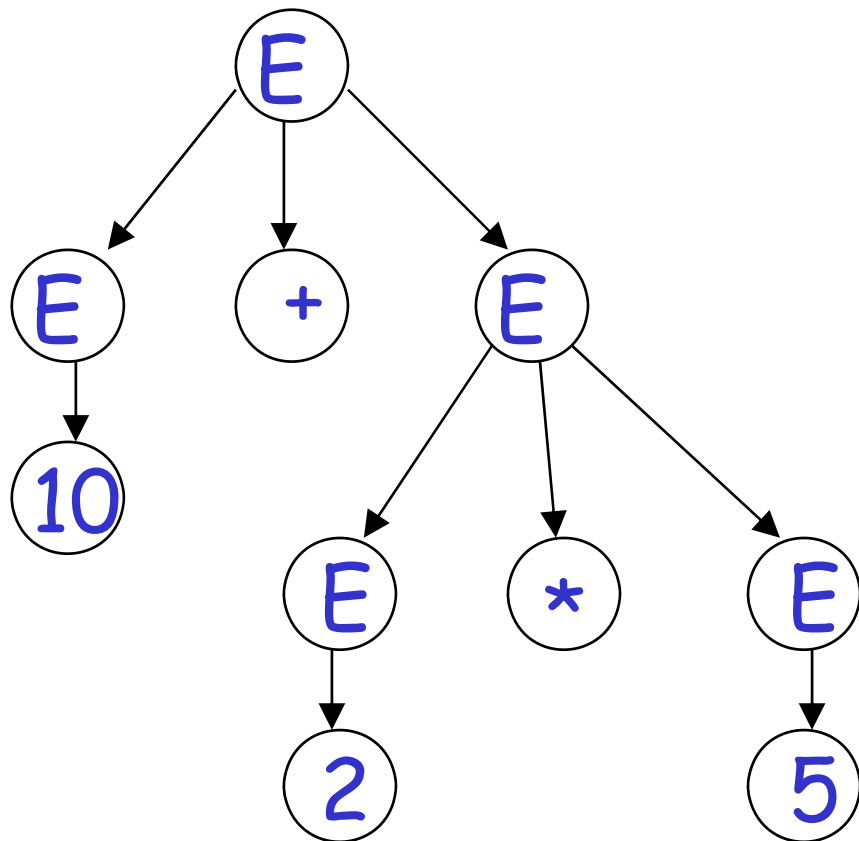
$E \Rightarrow E + E$
 $\Rightarrow E + E * E$
 $\Rightarrow 10 + E * E$
 $\Rightarrow 10 + 2 * E$
 $\Rightarrow 10 + 2 * 5$



derivation tree



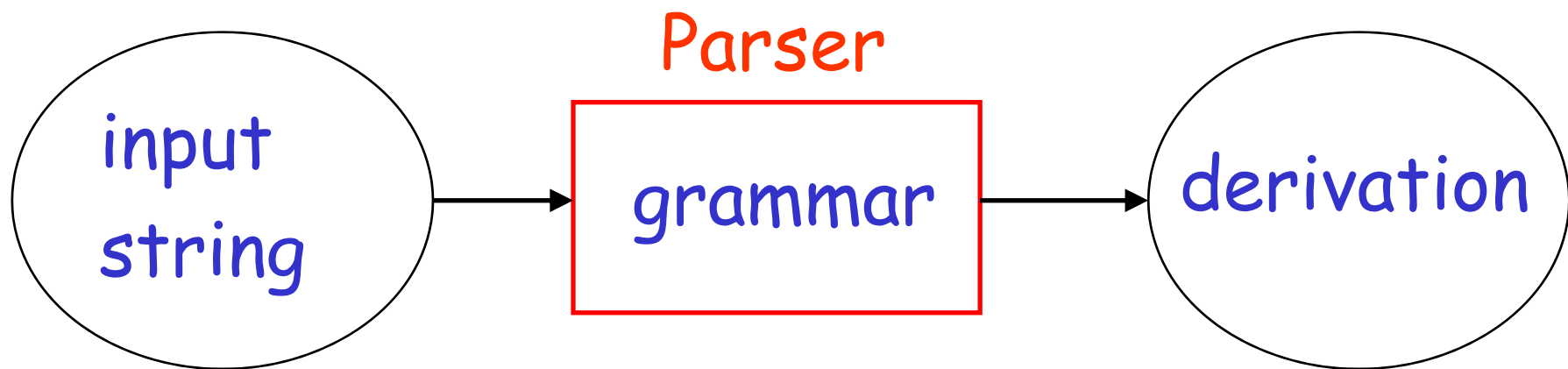
derivation tree



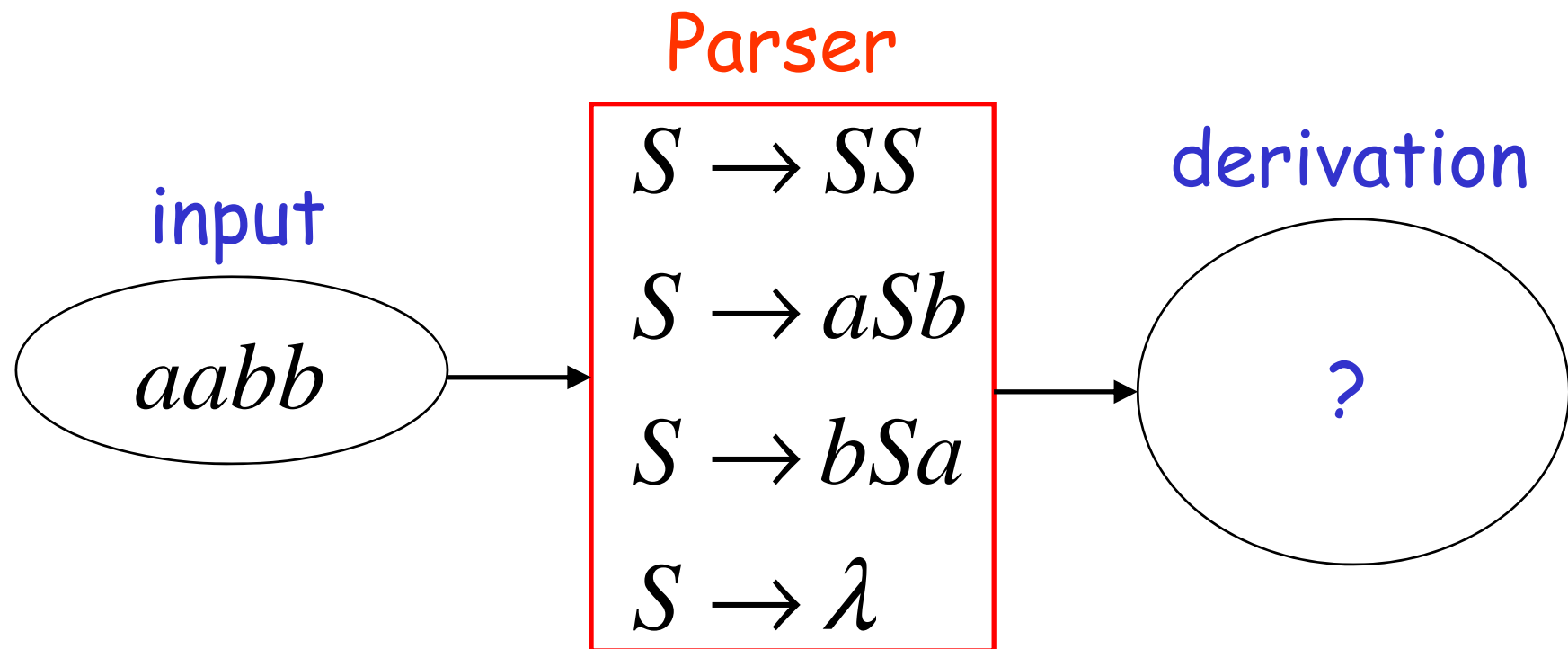
machine code

mult a, 2, 5
add b, 10, a

Parsing



Example:



Exhaustive Search

$$S \rightarrow SS \mid aSb \mid bSa \mid \lambda$$

Phase 1: $S \Rightarrow SS$ Find derivation of
 $S \Rightarrow aSb$ $aabb$
 $S \Rightarrow bSa$
 $S \Rightarrow \lambda$

All possible derivations of length 1

$$S \Rightarrow SS$$

aabb

$$S \Rightarrow aSb$$

~~$$S \Rightarrow bSa$$~~

~~$$S \Rightarrow \lambda$$~~

Phase 2 $S \rightarrow SS \mid aSb \mid bSa \mid \lambda$

$S \Rightarrow SS \Rightarrow SSS$

$S \Rightarrow SS \Rightarrow aSbS$

$aabb$

Phase 1

~~$S \Rightarrow SS \Rightarrow bSaS$~~

$S \Rightarrow SS$

$S \Rightarrow SS \Rightarrow S$

$S \Rightarrow aSb$

$S \Rightarrow aSb \Rightarrow aSSb$

$S \Rightarrow aSb \Rightarrow aaSbb$

~~$S \Rightarrow aSb \Rightarrow abSab$~~

~~$S \Rightarrow aSb \Rightarrow ab$~~

$$S \rightarrow SS \mid aSb \mid bSa \mid \lambda$$

Phase 2

$$S \Rightarrow SS \Rightarrow SSS$$

$$S \Rightarrow SS \Rightarrow aSbS$$

$$aabb$$

$$S \Rightarrow SS \Rightarrow S$$

$$S \Rightarrow aSb \Rightarrow aSSb$$

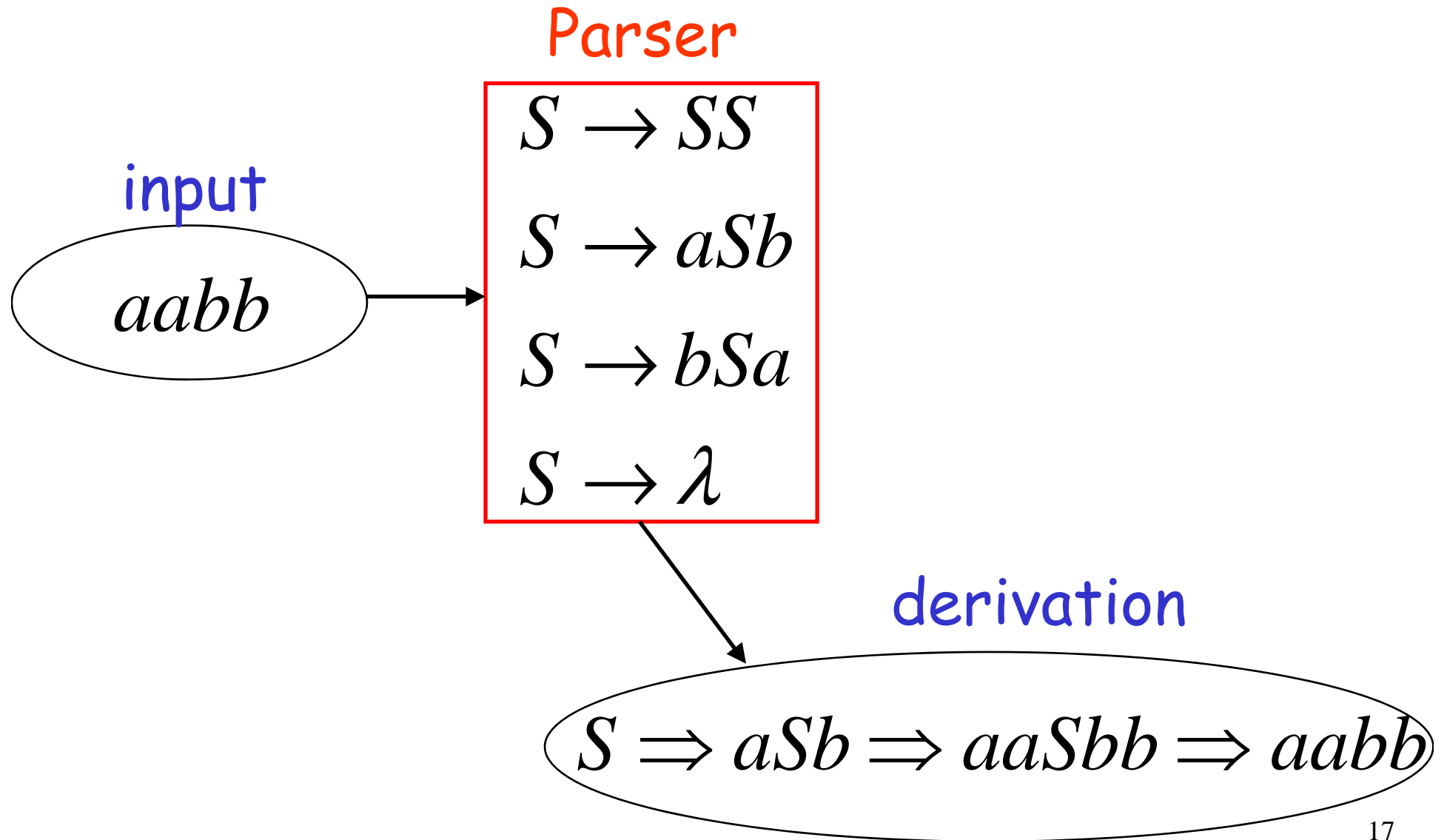
$$S \Rightarrow aSb \Rightarrow aaSbb$$

Phase 3



$$S \Rightarrow aSb \Rightarrow aaSbb \Rightarrow aabb$$

Final result of exhaustive search (top-down parsing)



Time complexity of exhaustive search

Suppose there are no productions of the form

$$A \rightarrow \lambda$$

$$A \rightarrow B$$

Number of phases for string w : $2^{|w|}$

For grammar with k rules

Time for phase 1: k

k possible derivations

Time for phase 2: k^2

k^2 possible derivations

Time for phase $2|w|$: $k^{2|w|}$

$k^{2|w|}$ possible derivations

Total time needed for string w :

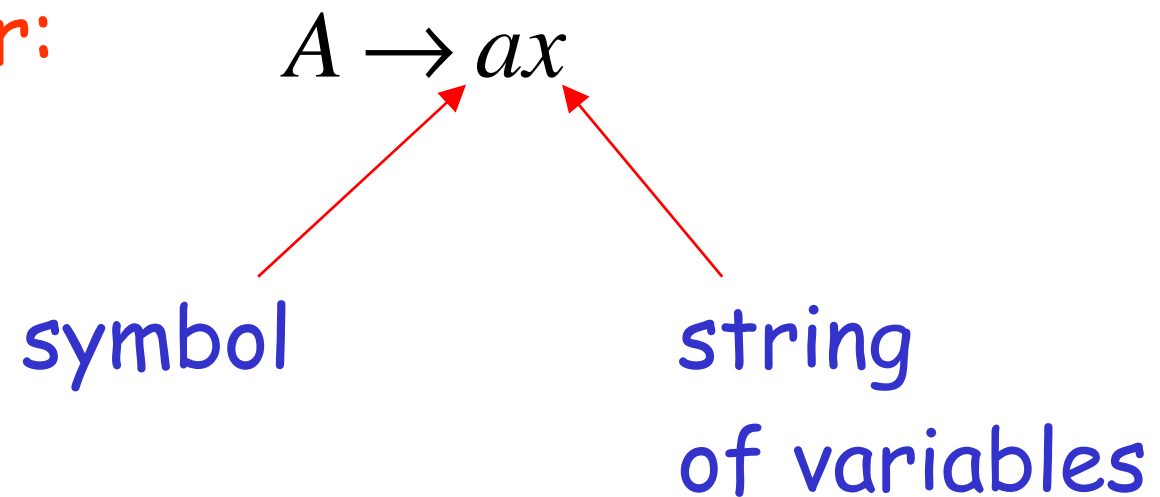
$$k + k^2 + \dots + k^{2|w|}$$

phase 1 phase 2 phase $2|w|$

Extremely bad!!!

There exist faster algorithms
for specialized grammars

S-grammar:



Pair (A, a) appears once

S-grammar example:

$$S \rightarrow aS$$

$$S \rightarrow bSS$$

$$S \rightarrow c$$

Each string has a unique derivation

$$S \Rightarrow aS \Rightarrow abSS \Rightarrow abcS \Rightarrow abcc$$

For S -grammars:

In the exhaustive search parsing
there is only one choice in each phase

Time for a phase: 1

Total time for parsing string w : $|w|$

For general context-free grammars:

There exists a parsing algorithm
that parses a string $|w|$
in time $|w|^3$